# An introduction to categorical probability theory

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It should be said: for someone trained in formal methods, the area of probability theory can be rather sloppy: everything is called 'P', types are hardly ever used, crucial ingredients (like distributions in expected values) are left implicit, basic notions (like conjugate prior) are introduced only via examples, calculation recipes and algorithms are regularly just given, without explanation, goal or justification, etc. This hurts, especially because there is so much beautiful mathematical structure around. (Jacobs [2019])

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Usually study the joint or marginal behaviour of X, Y, etc.

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- Kolmogorov-style conditioning is highly technical

# Practical implications





Implications:

- Lack of conceptual scalability that often requires hand-waving
- Difficult to interface with other mathematical theories
- Impediment to formal verification and automation
- A challenge pedagogically

Categorical probability reorganises the existing theory in a way that makes reasoning about higher-level concepts easy and intuitive

Theory becomes much more like a (high-level, expressive) programming language



#### PROBABILISTIC SYMMETRIES AND INVARIANT NEURAL NETWORKS

#### By Benjamin Bloem-Reddy $^1$ and Yee Whye $T\mathrm{EH}^2$

<sup>1</sup>Department of Statistics University of British Columbia benbr@stat.ubc.ca

<sup>2</sup>Department of Statistics University of Oxford y.w.teh@stats.ox.ac.uk Often it is desirable for a function  $f : \mathcal{X} \to \mathcal{Y}$  to be invariant to the action of a group  $\mathcal{G}$ 

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- $f: \mathcal{X} \to \mathcal{Y}$  makes some prediction about the population

Important question: for a given group  $\mathcal{G}$ , characterise the class of  $f: \mathcal{X} \to \mathcal{Y}$  such that

$$f(g \cdot x) = f(x)$$
 for all  $g \in \mathcal{G}$  and  $x \in \mathcal{X}$ 

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Aim is to characterise when Y is conditionally G-invariant in the sense that

$$\mathbb{P}(Y \in B \mid X \in A) = \mathbb{P}(Y \in B \mid X \in g \cdot A)$$

for all  $g \in \mathcal{G}$ ,  $A \in \Sigma_{\mathcal{X}}$  with  $\mathbb{P}(X \in A) > 0$ , and  $B \in \Sigma_{\mathcal{Y}}$ 

THEOREM 7. Let X and Y be random elements of Borel spaces X and Y, respectively, and G a compact group acting measurably on X. Assume that  $P_X$  is G-invariant, and pick a maximal invariant  $M : X \to S$ , with S another Borel space. Then  $P_{Y|X}$  is G-invariant if and only if there exists a measurable function  $f : [0,1] \times S \to Y$  such that

(14)  $(X,Y) \stackrel{\text{a.s.}}{=} (X, f(\eta, M(X)))$  with  $\eta \sim \text{Unif}[0,1]$  and  $\eta \perp \!\!\!\perp X$ .

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Here a maximal invariant is any measurable function M such that

$$M(x) = M(x') \Leftrightarrow x = g \cdot x'$$
 for some  $g \in \mathcal{G}$ 

(picture next slide)







Y

The proof of this is complex and uses highly technical ideas from advanced probability theory, e.g.

- Measurable cross section
- Normalised Haar measure
- Orbit law
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Also only applies when G is compact and X has a G-invariant marginal

### Why is this so hard to show? (E.g. compare deterministic case)

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Is it optimal to model a neural network in terms of random variables (X, Y)? And why must Law[X] be  $\mathcal{G}$ -invariant?

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Is it optimal to model a neural network in terms of random variables (X, Y)? And why must Law[X] be  $\mathcal{G}$ -invariant?

With the tools of categorical probability, we can not only generalise this result, but we can prove it in a way that maps directly onto our intuitions

# Categorical probability theory
A category (often) models a collection of entities that behave like functions:



Here X, Y, Z are objects and  $f : X \rightarrow Y$ ,  $g : Y \rightarrow X$  are arrows or morphisms

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Minimal structure:

- We can compose compatibly typed morphisms
- We have identity arrows

#### Definition

A category consists of a collection of objects and a collection of arrows

Each arrow f has a source and target object, denoted  $f : \mathcal{X} \rightarrow \mathcal{Y}$ 

There is a composition operation  $\circ$  on arrows such that

 $g \circ f : \mathcal{X} \to \mathcal{Z}$  whenever  $f : \mathcal{X} \to \mathcal{Y}$  and  $g : \mathcal{Y} \to \mathcal{Z}$  $h \circ (g \circ f) = (h \circ g) \circ f$  when f, g, h are appropriately typed

For every object  $\mathcal{X}$  there is an identity arrow  $id_{\mathcal{X}}: \mathcal{X} \to \mathcal{X}$  satisfying

 $f \circ \operatorname{id}_{\mathcal{X}} = f$  whenever  $f : \mathcal{X} \to \mathcal{Y}$  $\operatorname{id}_{\mathcal{X}} \circ g = g$  whenever  $g : \mathcal{Z} \to \mathcal{X}$ 

Philosophy: study structural properties extrinsically in terms of arrows

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- Set, the category of sets and functions
- Top, the category of topological spaces and continuous functions
- Meas, the category of measurable spaces and measurable functions

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- Stoch, the category of measurable spaces and Markov kernels
- A group  ${\mathcal G}$  can be viewed as a category (with a single object, and inverses)
- A poset can be viewed as a category (with a unique arrow  $x \to y$  iff  $x \leq y$ )

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Categories and functors themselves form a category...

#### Analogy



Denote by  $P\mathcal{X}$  the set of probability measures on  $\mathcal{X}$  (where  $\Sigma_{\mathcal{X}}$  implicit)

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• Equip  $P\mathcal{X}$  with the (initial)  $\sigma$ -algebra generated by the functions:

$$\mathsf{eval}_{\mathcal{A}}: \mathcal{PX} o [0,1] \qquad \mathsf{where} \ \mathcal{A} \in \Sigma_{\mathcal{X}} \ p \mapsto p(\mathcal{A})$$

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• For measurable  $f : \mathcal{X} \to \mathcal{Y}$ , define *Pf* by the pushforward, i.e.

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This reduces already the complexity of our original picture (since  $\mathbb{P} \in P\Omega$ )

Consider a measurable function  $k : \mathcal{X} \rightarrow P\mathcal{Y}$ 

Consider a measurable function  $k: \mathcal{X} \rightarrow \mathcal{PY}$ 

By definition of *P*:

- k(x)(-) is a probability measure for all  $x \in \mathcal{X}$
- $k(-)(B) = \operatorname{eval}_B \circ k$  is measurable for all  $B \in \Sigma_\mathcal{Y}$

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Hence k is a Markov kernel: can think of as  $k : \mathcal{X} \times \Sigma_{\mathcal{Y}} \to [0, 1]$  such that

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(Precisely: write  $k: \mathcal{X} \to P\mathcal{Y}$  as  $k: \mathcal{X} \to (\Sigma_{\mathcal{Y}} \mapsto [0, 1])$  and uncurry)

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Every "generalised generalised" function k : X → PPY can be canonically identified with E<sub>Y</sub> ∘ k : X → PY, where

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*P*,  $\delta_{\mathcal{Y}}$ , and  $E_{\mathcal{Y}}$  moreover satisfy coherence conditions and so give rise to a monad structure on **Meas** 

Given  $k : \mathcal{X} \to P\mathcal{Y}$  and  $\ell : \mathcal{Y} \to P\mathcal{Z}$ , define  $\ell \circ_{kl} k : \mathcal{X} \to P\mathcal{Z}$  via the following composition:

$$\mathcal{X} \xrightarrow{k} P\mathcal{Y} \xrightarrow{P\ell} PP\mathcal{Z} \xrightarrow{E_{\mathcal{Z}}} P\mathcal{Z}$$

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$$\mathcal{X} \stackrel{k}{\longrightarrow} \mathcal{P}\mathcal{Y} \stackrel{\mathcal{P}\ell}{\longrightarrow} \mathcal{P}\mathcal{P}\mathcal{Z} \stackrel{\mathcal{E}_{\mathcal{Z}}}{\longrightarrow} \mathcal{P}\mathcal{Z}$$

Can show this is the usual Chapman-Kolmogorov equation:

$$(\ell \circ_{\mathsf{kl}} k)(x)(A) = \int_{\mathcal{Y}} k(x)(\mathrm{d} y) \, \ell(y)(A) \qquad ext{where } A \in \Sigma_{\mathcal{Z}}$$

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Dirac maps  $\delta_{\mathcal{X}} : \mathcal{X} \to P\mathcal{X}, x \mapsto \text{Dirac}(x)$  behave like identities

#### This gives rise to the Kleisli category of Meas, known as Stoch:

	Meas	Stoch
Objects	Measurable spaces	Measurable spaces
Arrows	Measurable functions	Markov kernels
Composition	Composition of functions	Chapman-Kolmogorov
Identities	Identity functions	Dirac maps

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Interesting question: what Markov kernel corresponds to the measurable function  $id_{P\mathcal{Y}} : P\mathcal{Y} \to P\mathcal{Y}$ ?

$$\longleftrightarrow$$
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$$\mathsf{samp}_{\mathcal{Y}}: P\mathcal{Y} \to \mathcal{Y} \qquad \Longleftrightarrow \qquad \mathrm{id}_{P\mathcal{Y}}: P\mathcal{Y} \to P\mathcal{Y}$$

Here samp<sub> $\mathcal{Y}$ </sub>(p)(B) = p(B), i.e. samp<sub> $\mathcal{Y}$ </sub> draws a sample from its input

Stoch unifies and generalises the elements in our original picture:



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Although to some extent  $(\Omega, \Sigma_{\Omega})$  is redundant now . . .

# Return to case study

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# Conditional distributions/disintegrations

### Proposition

If  $\mathcal{Y}$  is standard Borel, then for any distribution p on  $\mathcal{X} \otimes \mathcal{Y}$ , there exists a Markov kernel  $k : \mathcal{X} \to \mathcal{Y}$  such that

$$p(A \times B) = \int_A \operatorname{proj}_{\mathcal{X}}(p)(\mathrm{d}x) \, k(x)(B)$$
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It is convenient to have a graphical way to denote this. Standard commutative diagrams get complex, but string diagrams work:



(Read from bottom to top)

# Informal usage

We use these informally all the time already, e.g. [Vaswani et al., 2017]:



Figure 1: The Transformer - model architecture.

### Definition

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For p = Law[X, Y], equivalent to conditional invariance in sense of Bloem-Reddy and Teh [2020] under their setup, i.e.  $\mathcal{G}$  is compact, Law[X]is  $\mathcal{G}$ -invariant,  $\mathcal{Y}$  standard Borel, and

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Makes sense more generally – could even start with k as the definition of a (probabilistic) neural network

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Explicitly, 
$$\Sigma_{\mathcal{X}/\sim} \coloneqq \{B \subseteq \mathcal{X}/\sim \mid q^{-1}(B) \in \Sigma_{\mathcal{X}}\}.$$

A measurable function  $g : \mathcal{X} \to \mathcal{Z}$  is  $\sim$ -invariant iff there exists a (necessarily unique) measurable function  $\tilde{g} : \mathcal{X}/\sim \to \mathcal{Z}$  such that  $\tilde{g} \circ q = g$ , i.e. the following diagram commutes:



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A very natural result in the context of category theory

Now take  $\mathcal{Z} = P\mathcal{Y}$  and interpret within **Stoch** 

### Corollary

A Markov kernel  $k : \mathcal{X} \to \mathcal{Y}$  is  $\sim$ -invariant iff there exists a Markov kernel  $\tilde{k} : \mathcal{X}/\sim \to \mathcal{Y}$  with



Now take  $\mathcal{Z} = P\mathcal{Y}$  and interpret within **Stoch** 



(Note that we are identifying q with its lifted version  $\delta_{\mathcal{X}/\sim} \circ q$ )

For any Markov kernel  $k : \mathcal{Z} \to \mathcal{Y}$  with  $\mathcal{Y}$  standard Borel, there exists a measurable function  $f : \mathcal{Z} \otimes [0,1] \rightarrow \mathcal{Y}$  such that



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Standard result (e.g. Lemma 3.22 of Kallenberg [2002])

If  $\mathcal{Y}$  is standard Borel, then  $\operatorname{Law}[X, Y]$  is conditionally  $\sim$ -invariant iff there exists a measurable function  $f : \mathcal{X}/\sim \otimes [0, 1] \rightarrow \mathcal{Y}$  such that

 $(X, Y) \stackrel{\mathrm{d}}{=} (X, f(q(X), \eta))$  where  $\eta \sim \mathrm{Uniform}(0, 1), \eta \perp X$ 

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**Proof**: writing p := Law[X, Y], conditional ~-invariance implies



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Conversely, right-hand side is conditionally  $\sim$ -invariant since q is.

THEOREM 7. Let X and Y be random elements of Borel spaces X and Y, respectively, and G a compact group acting measurably on X. Assume that  $P_X$  is G-invariant, and pick a maximal invariant  $M : X \to S$ , with S another Borel space. Then  $P_{Y|X}$  is G-invariant if and only if there exists a measurable function  $f : [0,1] \times S \to Y$  such that

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Not quite done:

$$(X, Y) \stackrel{\mathrm{d}}{=} (X, f(q(X), \eta)) \qquad \Rightarrow \qquad Y \stackrel{\mathrm{a.s.}}{=} f(q(X), \eta)$$

Choose  $h:\mathcal{X}\otimes\mathcal{Y}\otimes[0,1]
ightarrow [0,1]$  such that



Existence of *h* follows by disintegrating right-hand side along  $\mathcal{X} \times \mathcal{Y}$  and then applying noise outsourcing result

# Completing the proof

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#### Theorem (Our version)

If  $\mathcal{Y}$  is Borel, then  $\operatorname{Law}[X, Y]$  is conditionally  $\sim$ -invariant iff there exists a measurable function  $f : \mathcal{X}/\sim \otimes [0, 1] \to \mathcal{Y}$  such that

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Possibly better to express entirely via Markov kernels



Categorical probability offers a high-level perspective on the classical theory that makes abstraction easier and helps theory follow intuition

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The programming language has been (increasingly) written – now is time for practitioners to write new software

## Equivariant stochastic neural networks in Markov categories

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